

Computational algebraic geometry

Learning coefficients via symbolic and numerical methods

Anton Leykin

Georgia Tech

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Ideals \leftrightarrow varieties

- Let k be a field (\mathbb{R} or \mathbb{C}).
- Ideal** in $R = k[x_1, \dots, x_n]$ generated by f_1, \dots, f_r is

$$I = \langle f_1, \dots, f_r \rangle = \sum_{i=1}^r g_i f_i \subset R$$

- Variety** defined by I is

$$V = \mathbb{V}(I) = \{p \in k^n : \forall f \in I, f(p) = 0\}$$

- The **vanishing ideal** of a variety V is

$$I = \mathbb{I}(V) = \{f \in R : f(V) = 0\}$$

- The **radical** of I is

$$\sqrt{I} = \{g : \exists m, g^m \in I\} \subset R, \quad \mathbb{V}(I) = \mathbb{V}(\sqrt{I})$$

I is a radical ideal if $I = \sqrt{I}$

- Hilbert Nullstellensatz:** if $k = \mathbb{C}$ then $\mathbb{I}(\mathbb{V}(J)) = \sqrt{J}$ for all $J \subset R$.

Gröbner bases, invariants

- A **monomial order** $>$ is given by weights $w_1, w_2, \dots \in \mathbb{R}_{\geq 0}^n$:

$$x^\alpha > x^\beta \Leftrightarrow w_1 \cdot \alpha > w_1 \cdot \beta \text{ or} \\ w_2 \cdot \alpha > w_2 \cdot \beta \text{ or} \\ \dots$$

such that $>$ is total.

Example: $>_{lex}$ is given by $w_1 = e_1, w_2 = e_2, \dots$

- The initial term of $f \in R$ is the term (coefficient·monomial) with the largest monomial occurring with a nonzero coefficient:

e.g. for $f = 3x^2y + 4xy^7 + 2$ it is $\text{in}_{lex}(f) = 3x^2y$.

- A **Gröbner basis** of I is a set of generators G of I such that

$$\langle \text{in}(G) \rangle = \langle \text{in}(I) \rangle.$$

- **Invariants** such as dimension, degree, Hilbert polynomial can be computed using Gröbner bases.

Computational (symbolic) algebraic geometry

- Cox, Little, O'Shea, [Ideals, varieties, and algorithms](#)
- Cox, Little, O'Shea, [Using algebraic geometry](#)
- Greuel, Pfister, [A Singular Introduction to Commutative Algebra](#)

Software:

- [Singular](#) (can resolve singularities)
- [Macaulay2](#) (can do everything else)
- [Magma](#) (Australia), [CoCoA](#) (Italy), [risa/asir](#) (Japan), ... (other specialized software)
- Sage, Maple, Mathematica, ... (more general software)

Polynomial homotopy continuation

- **Target** system: n equations in n variables,

$$F(\mathbf{x}) = (f_1(\mathbf{x}), \dots, f_n(\mathbf{x})) = \mathbf{0},$$

where $f_i \in R = \mathbb{C}[\mathbf{x}] = \mathbb{C}[x_1, \dots, x_n]$ for $i = 1, \dots, n$.

- **Start** system: n equations in n variables:

$$G(\mathbf{x}) = (g_1(\mathbf{x}), \dots, g_n(\mathbf{x})) = \mathbf{0},$$

such that it is easy to solve.

- **Homotopy**: for $\gamma \in \mathbb{C} \setminus \{0\}$ consider

$$H(\mathbf{x}, t) = (1 - t)G(\mathbf{x}) + \gamma t F(\mathbf{x}), \quad t \in [0, 1].$$

Example

target

$$f_1 = x_1^4 x_2 + 5x_1^2 x_2^3 + x_1^3 - 4$$

$$f_2 = x_1^2 - x_1 x_2 + x_2 - 8$$

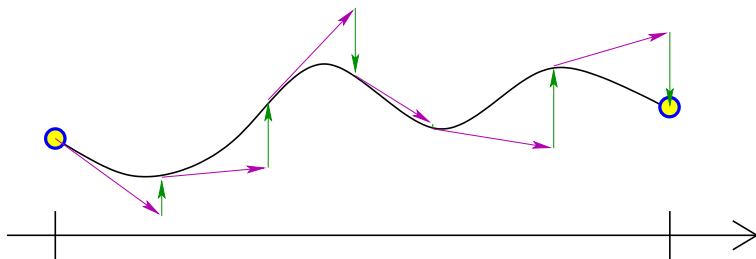
start

$$g_1 = x_1^5 - 1$$

$$g_2 = x_2^2 - 1$$

Start solutions \rightarrow target solutions:

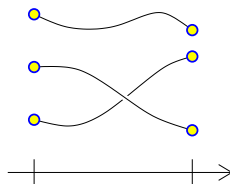
$$H(\mathbf{x}, t) = 0 \text{ implies } \frac{d\mathbf{x}}{dt} = - \left(\frac{\partial H}{\partial \mathbf{x}} \right)^{-1} \frac{\partial H}{\partial t}.$$



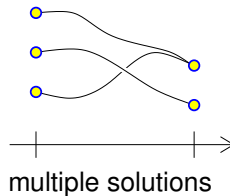
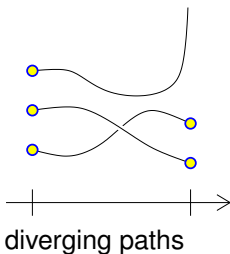
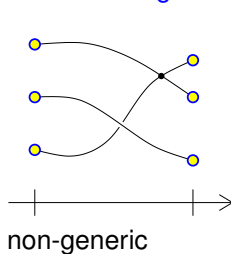
Global picture

Optimal homotopy:

- the continuation paths are **regular**;
- the homotopy establishes a bijection between the start and target solutions.



Possible **singular** scenarios:

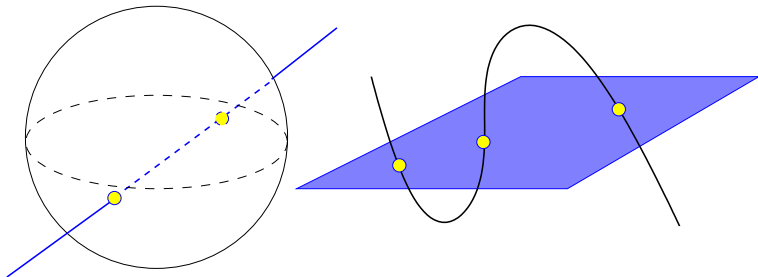


Higher-dimensional solution sets

- Let $I = (f_1, \dots, f_N)$ be an ideal of $\mathbb{C}[x_1, \dots, x_n]$.
- **Goal:** Understand the variety

$$X = \mathbb{V}(I) = \{\mathbf{x} \in \mathbb{C}^n \mid \forall f \in I, f(\mathbf{x}) = 0\}.$$

- A **witness set** for an equidimensional component Y of X :
 - a generic “slicing” plane L with $\dim L = \text{codim } Y$
 - witness points $w_{Y,L} = Y \cap L$
 - (generators of I)



Numerical algebraic geometry

- Sommese, Verschelde, and Wampler, [Introduction to Numerical AG](#) (2005)
- Sommese and Wampler, [The numerical solution of systems of polynomials](#) (2005)

Software:

- [PHCpack](#) (Verschelde);
- [HOM4PS](#) (group of T.Y.Li);
- [Bertini](#) (group of Sommese);
- [NAG4M2](#): Numerical Algebraic Geometry for Macaulay2 (L.).

and more, e.g.: [Maple](#)'s `ROOTFINDING[HOMOTOPY]`.

Analysis

- For a polynomial f , the function $|f|^s$ is locally integrable $\operatorname{Re} s > 0$.
- Hence, $|f|^s$ is a generalized function defined on $\{s : \operatorname{Re} s > 0\} \subset \mathbb{C}$: for a distribution $\varphi \in C_c^\infty$,

$$|f|^s(\varphi) = \int |f(x)|^s \varphi(x) dx$$

- **Gelfand [1957]**: Does it extend to a meromorphic function on \mathbb{C} ?
- **I. N. Bernstein [1968]**: Yes. The poles are contained in a finite number of arithmetic progressions.
- Key ingredients: **resolution of singularities** and being able to write a **functional equation**

$$b(s)f^s = P \cdot f^{s+1}$$

when f is a monomial. (Here: $b(s)$ is a univariate polynomial and P is a linear differential operator with coefficients in $\mathbb{C}[\mathbf{x}, s]$.)

Invariants in singularity theory

Definition (Multiplier ideal for $\mathbf{f} = (f_1, \dots, f_r)$)

$$\mathcal{J}(\mathbf{f}^c) = \left\{ h \in \mathbb{C}[\mathbf{x}] : \frac{|h|^2}{(\sum |f_i|^2)^c} \text{ is locally integrable} \right\}.$$

For $r = 1$, it is the ideal of h , that make $\frac{|h|}{|f_1|^c}$ locally integrable.

- Algebrao-geometric definition: via log-canonical resolutions.
- **Jumping coefficients of \mathbf{f}** : rational numbers

$$0 = \xi_0 < \xi_1 < \xi_2 < \dots$$

such that $\mathcal{J}(\mathbf{f}^c)$ is constant exactly for $c \in [\xi_i, \xi_{i+1})$.

- ξ_1 is called the **log-canonical threshold**.
- These invariants measure singularities of the corresponding variety; in particular, they depend only on the ideal $\langle \mathbf{f} \rangle$.

Weyl algebra

- Let K be a field of characteristic zero. (Think: $K = \mathbb{C}$)
- Affine space: $X = K^n$.
- **Weyl algebra**: an associative algebra

$$D_X = K\langle \mathbf{x}, \boldsymbol{\partial} \rangle = K\langle x_1, \dots, x_n, \partial_1, \dots, \partial_n \rangle$$

where $[\partial_i, x_i] = \partial_i x_i - x_i \partial_i = 1$ and all other pairs of generators commute.

- D_X is isomorphic to the algebra of **linear differential operators** with polynomial coefficients.
- Every element has the **normal form**

$$Q = \sum_{\alpha, \beta \in \mathbb{Z}^n} c_{\alpha\beta} \mathbf{x}^\alpha \boldsymbol{\partial}^\beta,$$

where finitely many of $c_{\alpha\beta} \in K$ are nonzero.

D -modules

- D_X is simple: only trivial two-sided ideals.
- We consider only **left** ideals and **left** D_X -modules.
- Examples of D -modules: $K[x]$, $K[[x]]$, $C^\infty(X)$.
- **Software:**
 - kan/sml (Takayama)
 - risa/asir (Noro)
 - dmod.lib, Singular (Levandovsky et al.)
 - D-modules, Macaulay2 (L., Tsai)

Gröbner bases

- D_X is **Gröbner-friendly**: D_X is an algebra of **solvable type**.
- Gröbner bases can be computed with respect to any w -compatible monomial order, where $w = (w_x, w_\partial) \in \mathbb{R}^{2n}$ satisfies $w_x + w_\partial \geq 0$ componentwise.
- The **Bernstein-Sato** polynomial $b_f(s) \neq 0$ is the monic polynomial $b(s)$ of the minimal degree satisfying

$$b(s)f^s = P \cdot f^{s+1} = Pf \cdot f^s \in D_X[s]f^s,$$

where $P \in D_X[s]$ and $D_X[s]f^s$ is a cyclic $D_X[s]$ -module generated by f^s .

- One can compute $b_f(s)$ via Gröbner bases.

Connection between $\text{lct}(f)$ and $b_f(s)$

- Assume $k = \mathbb{C}$ then
 - The log-canonical threshold c_0 is the lowest root of $b_f(-s)$.
 - Every jumping coefficient $c \in [c_0, c_0 + 1)$ is a root of $b_f(-s)$.
- For $k = \mathbb{R}$ there are examples where $\text{rlct}(f)$ is not a jumping coefficient of f .
However, $\text{rlct}(f) \bmod 1$ equals some root of $b_f(-s)$.
- M2 packages where lct is computed:
 - `Dmodules` (general case)
 - `MonomialMultiplierIdeals` (monomial ideals)
 - `HyperplaneArrangements` (central hyperplane arrangements)

Improper integrals

To find $c = \text{rlct}(f)$ look for the pole $s = -c$ closest to 0 of

$$|f|^s(\varphi) = \int_{x \in \mathbb{R}^n} |f(x)|^s \varphi(x) dx.$$

Idea:

- Pick φ supported on a neighborhood of a singularity of the variety $\{f = 0\}$ and evaluate $|f|^s(\varphi)$ **numerically**.
- If $|f|^{-c_1}(\varphi) < \infty$ and $|f|^{-c_2}(\varphi) = \infty$ then $c_1 < \text{rlct}(f) \leq c_2$.
- We can determine possible values of $\text{rlct}(f)$ by looking at $b_f(s)$?

Questions

- Can we say something about the denominator of $\text{rlct}(f)$?
- How to compute an improper integral numerically?
How to determine it equals ∞ ?
- Can we use statistics to estimate rlct ?
- Is there a hybrid strategy (e.g., “partial” resolution + b-functions)?